Consistency algorithms

Chapter 3

Consistency methods

- Approximation of inference:
  - Arc, path and i-consistecy
- Methods that transform the original network into a tighter and tighter representations

Arc-consistency

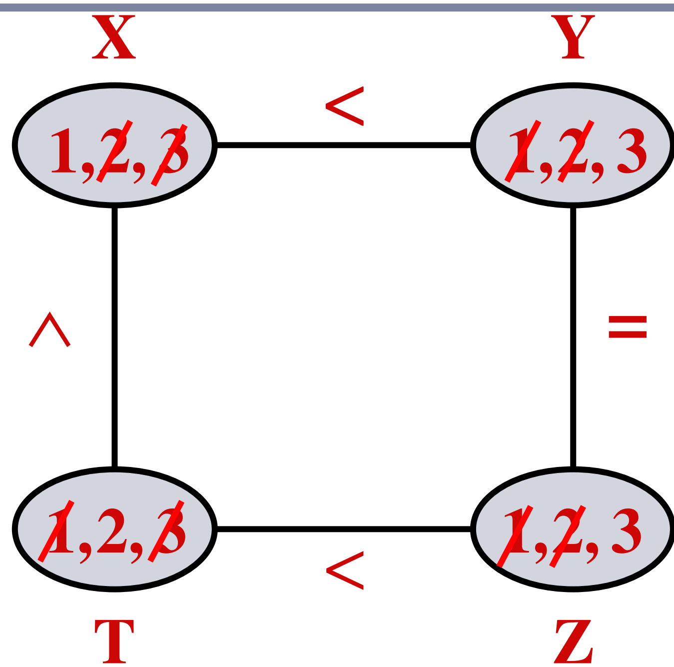
$1 \le X, Y, Z, T \le 3$

$X < Y$

$Y = Z$

$T < Z$

$X \le T$



Arc-consistency

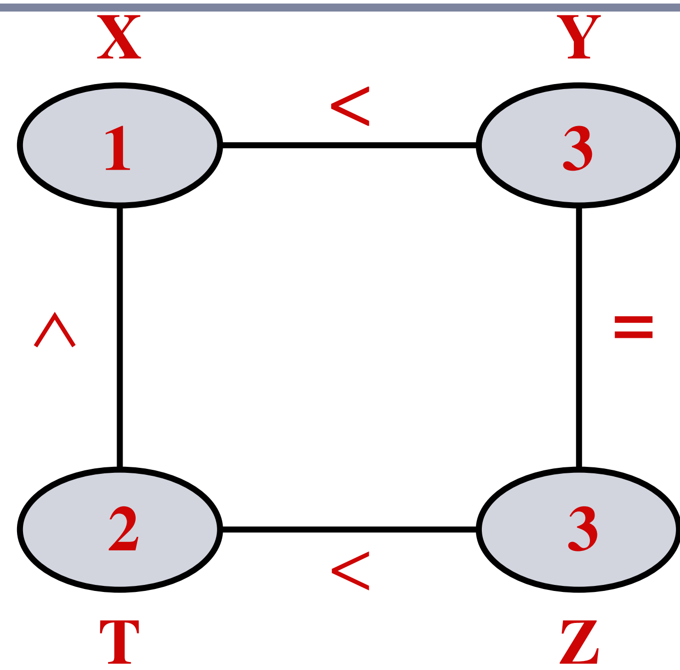
$1 \le X, Y, Z, T \le 3$

$X < Y$

$Y = Z$

$T < Z$

$X \le T$



Arc-consistency

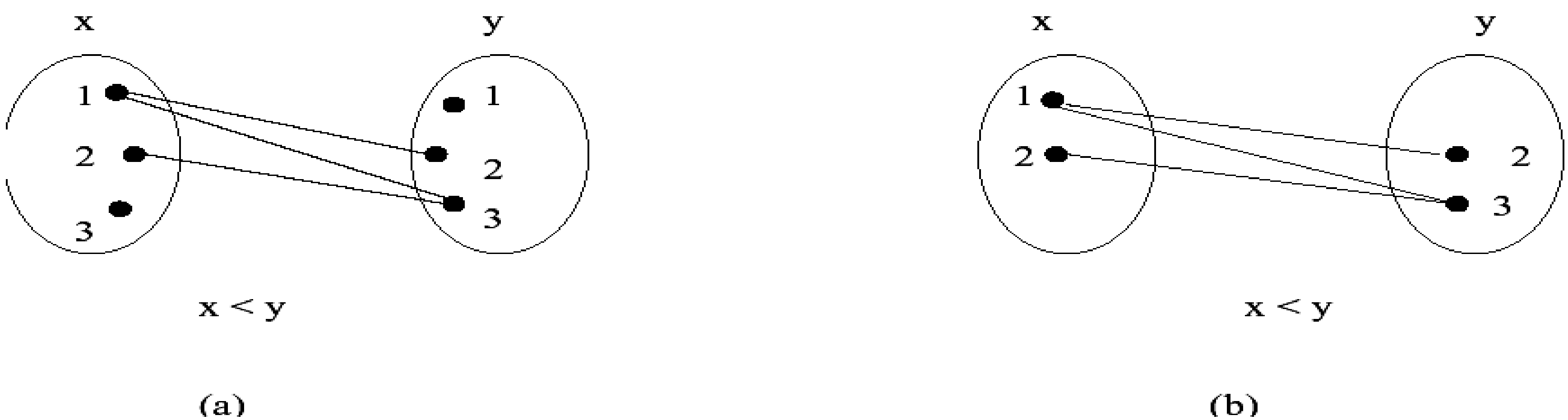


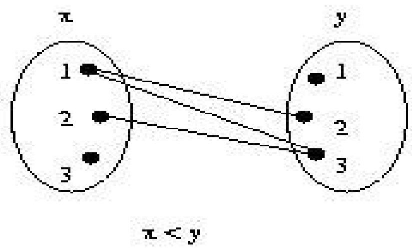
Figure 3.1: A matching diagram describing the arc-consistency of two variables x and y. In (a) the variables are not arc-consistent. In (b) the domains have been reduced, and the variables are now arc-consistent.

Definition 3.2.2 (arc-consistency) *Given a constraint network* $\mathcal{R} = (\mathcal{X}, \mathcal{D}, \mathcal{C})$, *with* $R_{ij} \in C$, *a variable* x_i *is* arc-consistent *relative to* x_j *if and only if for every value* $a_i \in D_i$ *there exists a value* $a_j \in D_j$ *such that* $(a_i, a_j) \in R_{ij}$. *The subnetwork (alternatively, the arc) defined by* $\{x_i, x_j\}$ *is arc-consistent if and only if* x_i *is arc-consistent relative to* x_j *and* x_j *is arc-consistent relative to* x_i. *A network of constraints is called* arc-consistent *iff all of its arcs (e.g., subnetworks of size 2) are arc-consistent.*

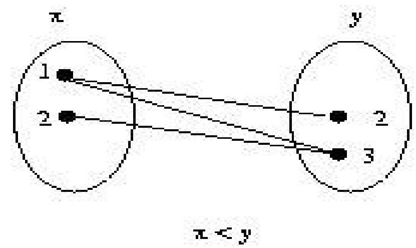
Arc-consistency

Only domain constraints are recorded: $R_A \leftarrow \prod_A R_{AB} \bowtie D_B$

Example: $R_X = \{1,2,3\}, R_Y = \{1,2,3\}$, constriant $X < Y$ reduces domain of X to $R_X = \{1,2\}$.



(a)



(b)

Revise for arc-consistency

REVISE($(x_i), x_j$)

input: a subnetwork defined by two variables $X = \{x_i, x_j\}$, a distinguished variable x_i, domains: D_i and D_j, and constraint R_{ij}

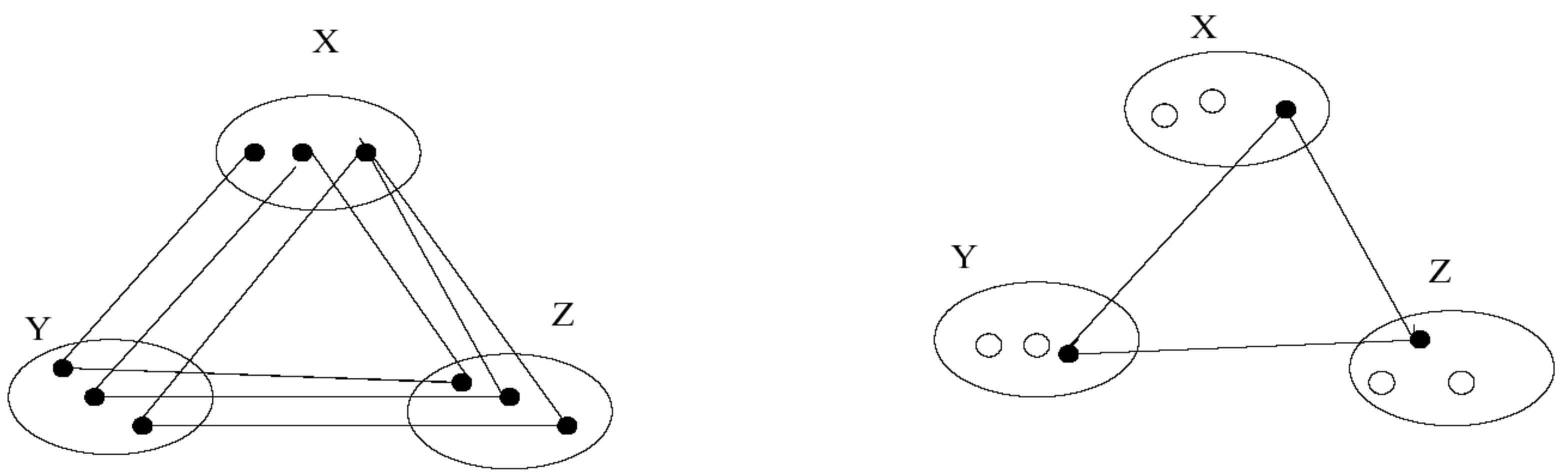
output: D_i, such that, x_i arc-consistent relative to x_j

1. **for** each $a_i \in D_i$
2. **if** there is no $a_j \in D_j$ such that $(a_i, a_j) \in R_{ij}$
3. **then** delete a_i from D_i
4. **endif**
5. **endfor**

Figure 3.2: The Revise procedure

$$D_i \leftarrow D_i \cap \pi_i (R_{ij} \otimes D_j)$$

Figure 3.3: (a) Matching diagram describing a network of constraints that is not arc-consistent (b) An arc-consistent equivalent network.



(a) (b)

AC-1

AC-1($\mathcal{R}$)

input: a network of constraints $\mathcal{R} = (X, D, C)$
output: $\mathcal{R}'$ which is the loosest arc-consistent network equivalent to $\mathcal{R}$

1. **repeat**
2. **for** every pair $\{x_i, x_j\}$ that participates in a constraint
3. Revise($(x_i), x_j$) (or $D_i \leftarrow D_i \cap \pi_i(R_{ij} \bowtie D_j)$)
4. Revise($(x_j), x_i$) (or $D_j \leftarrow D_j \cap \pi_j(R_{ij} \bowtie D_i)$)
5. **endfor**
6. **until** no domain is changed

Figure 3.4: Arc-consistency-1 (AC-1)

- Complexity (Mackworth and Freuder, 1986): $O(enk^3)$
- e = number of arcs, n variables,k values
- (ek^2, each loop, nk number of loops), best-case = ek,
- Arc-consistency is: $\Omega(ek^2)$

AC-3

AC-3($\mathcal{R}$)

input: a network of constraints $\mathcal{R} = (X, D, C)$

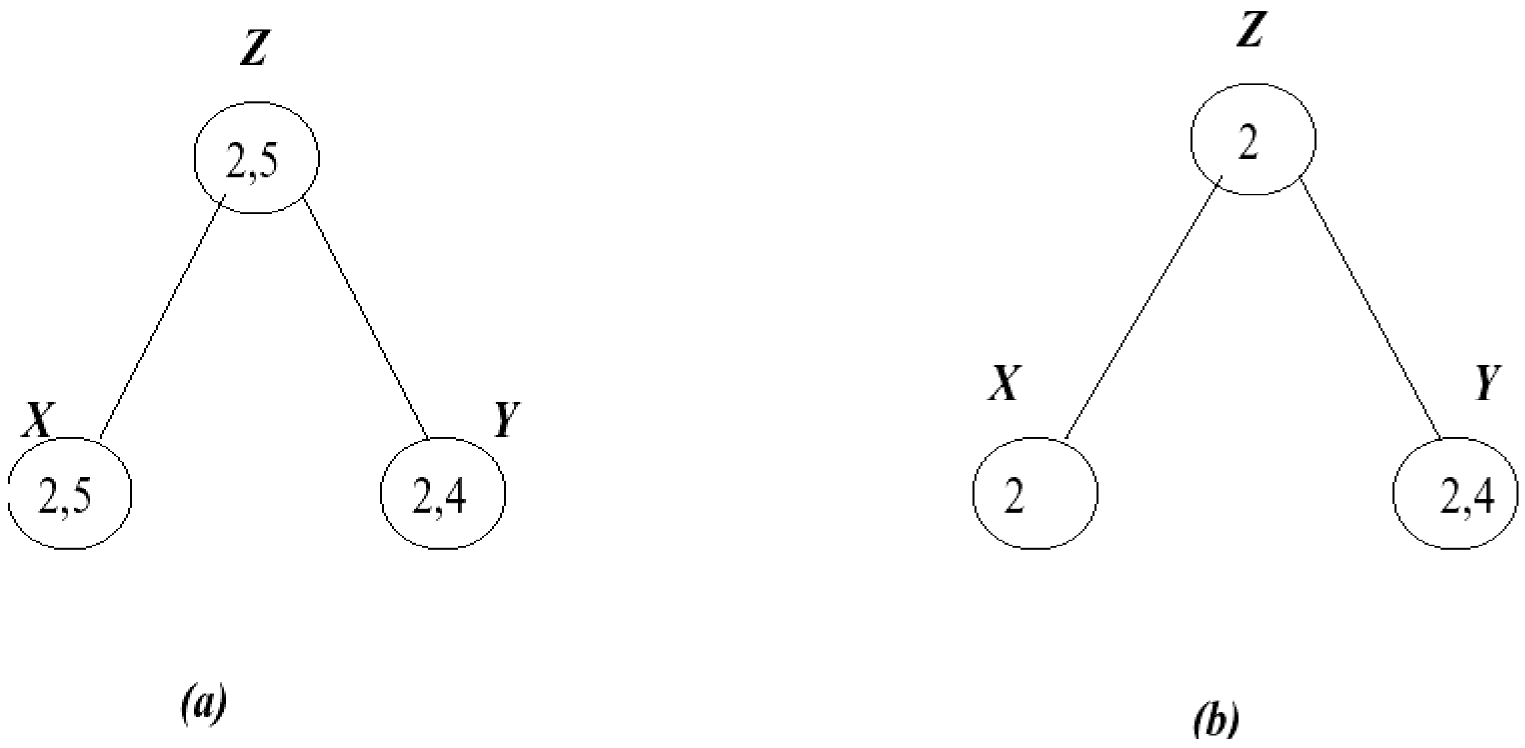
output: $\mathcal{R}'$ which is the largest arc-consistent network equivalent to $\mathcal{R}$

1. **for** every pair $\{x_i, x_j\}$ that participates in a constraint $R_{ij} \in \mathcal{R}$
2. $\quad$ $queue \leftarrow queue \cup \{(x_i, x_j), (x_j, x_i)\}$
3. **endfor**
4. **while** $queue \neq \{\}$
5. $\quad$ select and delete (x_i, x_j) from *queue*
6. $\quad$ $Revise((x_i), x_j)$
7. $\quad$ **if** $Revise((x_i), x_j)$ causes a change in D_i
8. $\quad\quad$ **then** $queue \leftarrow queue \cup \{(x_k, x_i), i \neq k\}$
9. $\quad$ **endif**
10. **endwhile**

Figure 3.5: Arc-consistency-3 (AC-3)

- Complexity: $O(ek^3)$
- Best case O(ek), since each arc may be processed in O(2k)

Example: A three variable network, with two constraints: z divides x and z divides y (a) before and (b) after AC-3 is applied.



AC-4

AC-4($\mathcal{R}$)
input: a network of constraints $\mathcal{R}$
output: An arc-consistent network equivalent to $\mathcal{R}$

Initialization: $M \leftarrow \emptyset$,
    initialize $S_{(x_i,c_i)}$, $counter(i, a_i, j)$ for all R_{ij}
    for all counters
        if $counter(x_i, a_i, x_j) = 0$ (if $< x_i, a_i >$ is unsupported by x_j)
            then add $< x_i, a_i >$ to $LIST$
        endif
    endfor
while $LIST$ is not empty
    choose $< x_i, a_i >$ from LIST, remove it, and add it to M
    for each $< x_j, a_j >$ in $S_{(x_i,a_i)}$
        decrement $counter(x_j, a_j, x_i)$
        if $counter(x_j, a_j, x_i) = 0$
            then add $< x_j, a_j >$ to $LIST$
        endif
    endfor
endwhile

Figure 3.7: Arc-consistency-4 (AC-4)

- Complexity: $O(ek^2)$
- (Counter is the number of supports to ai in xi from xj. S_(xi,ai) is the set of pairs that (xi,ai) supports)

Distributed arc-consistency (Constraint propagation)

- Implement AC-1 distributedly.

$$D_i \leftarrow D_i \cap \pi_i(R_{ij} \otimes D_j)$$

- Node x_j sends the message to node x_i

$$h_i^j \leftarrow \pi_i(R_{ij} \otimes D_j)$$

- Node x_i updates its domain:

$$D_i \leftarrow D_i \cap \pi_i(R_{ij} \otimes D_j) =$$

$$D_i \leftarrow D_i \cap h_i^j$$

- Messages can be sent asynchronously or scheduled in a topological order

Is arc-consistency enough?

- Example: a triangle graph-coloring with 2 values.
  - Is it arc-consistent?
  - Is it consistent?
- It is not path, or 3-consistent.

Path-consistency

Definition 3.3.2 (Path-consistency) *Given a constraint network $\mathcal{R} = (X, D, C)$, a two variable set $\{x_i, x_j\}$ is path-consistent relative to variable x_k if and only if for every consistent assignment $(< x_i, a_i >, < x_j, a_j >)$ there is a value $a_k \in D_k$ s.t. the assignment $(< x_i, a_i >, < x_k, a_k >)$ is consistent and $(< x_k, a_k >, < x_j, a_j >)$ is consistent. Alternatively, a binary constraint R_{ij} is path-consistent relative to x_k iff for every pair $(a_i, a_j), \in R_{ij}$, where a_i and a_j are from their respective domains, there is a value $a_k \in D_k$ s.t. $(a_i, a_k) \in R_{ik}$ and $(a_k, a_j) \in R_{kj}$. A subnetwork over three variables $\{x_i, x_j, x_k\}$ is path-consistent iff for any permutation of (i, j, k), R_{ij} is path consistent relative to x_k. A network is path-consistent iff for every R_{ij} (including universal binary relations) and for every $k \neq i, j$ R_{ij} is path-consistent relative to x_k.*

Path-consistency

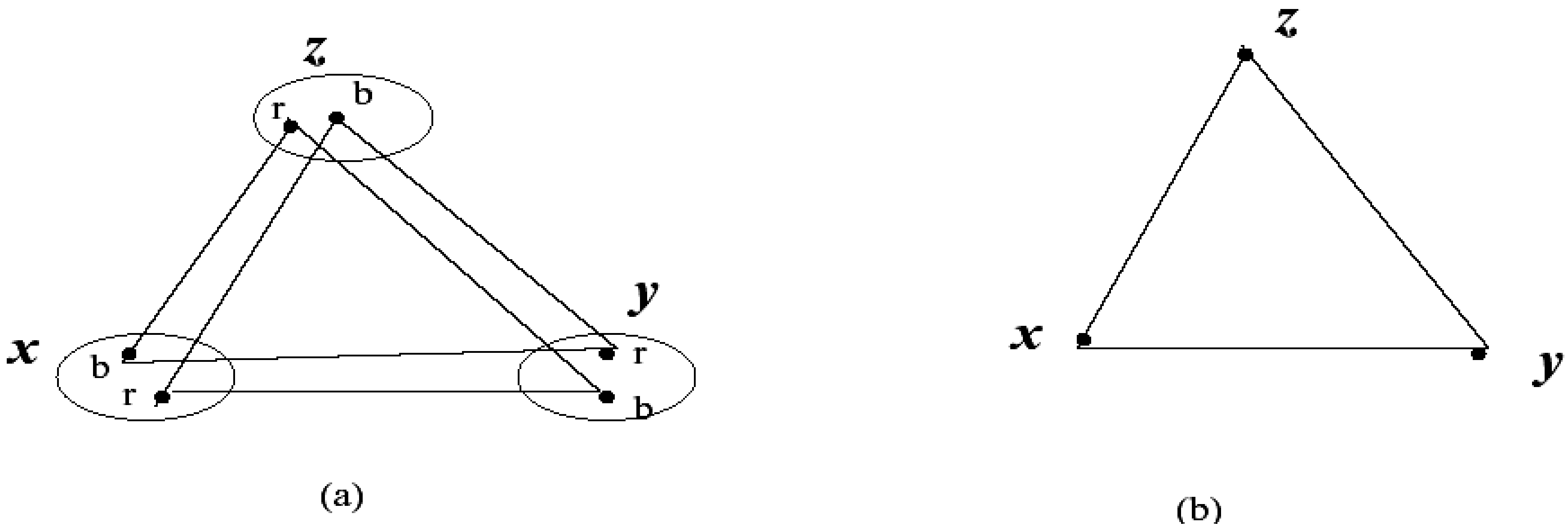


Figure 3.8: (a) The matching diagram of a 2-value graph coloring problem. (b) Graphical picture of path-consistency using the matching diagram.

Revise-3

REVISE-3$((x,y),z)$

input: a three-variable subnetwork over (x,y,z), R_{xy}, R_{yz}, R_{xz}.
output: revised R_{xy} path-consistent with z.

1. **for** each pair $(a,b) \in R_{xy}$
2.   **if** no value $c \in D_z$ exists such that $(a,c) \in R_{xz}$ and $(b,c) \in R_{yz}$
3.   **then** delete (a,b) from R_{xy}.
4.   **endif**
5. **endfor**

Figure 3.9: Revise-3

$$R_{ij} \leftarrow R_{ij} \cap \pi_{ij}(R_{ik} \otimes D_k \otimes R_{kj})$$

- Complexity: O(k^3)
- Best-case: O(t)
- Worst-case O(tk)

PC-1

PC-1($\mathcal{R}$)

input: a network $\mathcal{R} = (X, D, C)$.
output: a path consistent network equivalent to $\mathcal{R}$.

1. **repeat**
2.   **for** $k \leftarrow 1$ to n
3.   **for** $i, j \leftarrow 1$ to n
4. $R_{ij} \leftarrow R_{ij} \cap \pi_{ij}(R_{ik} \bowtie D_k \bowtie R_{kj}) /^* \ (Revise - 3((i,j),k))$
5.   **endfor**
6.   **endfor**
7. **until** no constraint is changed.

Figure 3.10: Path-consistency-1 (PC-1)

- Complexity: $O(n^5 k^5)$
- O(n^3) triplets, each take O(k^3) steps → O(n^3 k^3)
- Max number of loops: O(n^2 k^2) .

PC-2

PC-3($\mathcal{R}$)

input: a network $\mathcal{R} = (X, D, C)$.

output: $\mathcal{R}'$ a path consistent network equivalent to $\mathcal{R}$.

1. $Q \leftarrow \{(i,k,j) \mid 1 \le i < j \le n, 1 \le k \le n, k \ne i, k \ne j\}$
2. **while** Q is not empty
3. select and delete a 3-tuple (i,k,j) from Q
4. $R_{ij} \leftarrow R_{ij} \cap \pi_{ij}(R_{ik} \bowtie D_k \bowtie R_{kj})$ /* (Revise-3($(i,j),k$))
5. **if** R_{ij} changed then
6. $Q \leftarrow Q \cup \{(l,i,j)(l,j,i) \mid 1 \le l \le n, l \ne i, l \ne j\}$
7. **endwhile**

Figure 3.11: Path-consistency-3 (PC-3)

- Complexity: $O(n^3 k^5)$
- Optimal PC-4: $O(n^3 k^3)$
- (each pair deleted may add: 2n-1 triplets, number of pairs: O(n^2 k^2) → size of Q is O(n^3 k^2), processing is O(k^3))

Example: before and after path-consistency

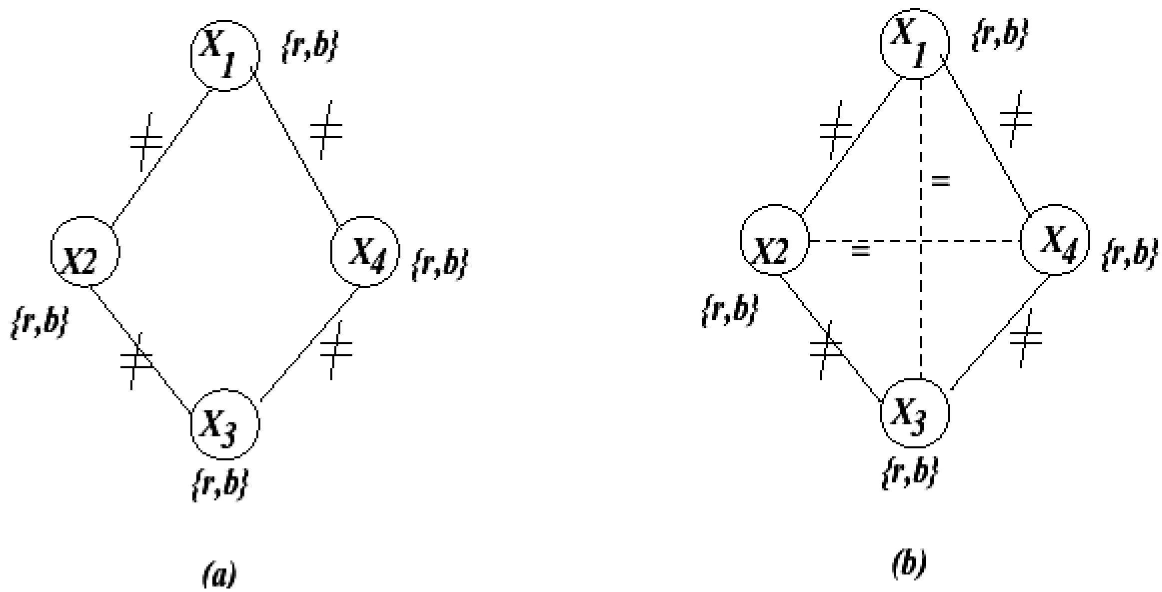


Figure 3.12: A graph-coloring graph (a) before path-consistency (b) after path-consistency

- PC-1 requires 2 processings of each arc while PC-2 may not
- Can we do path-consistency distributedly?

I-consistency

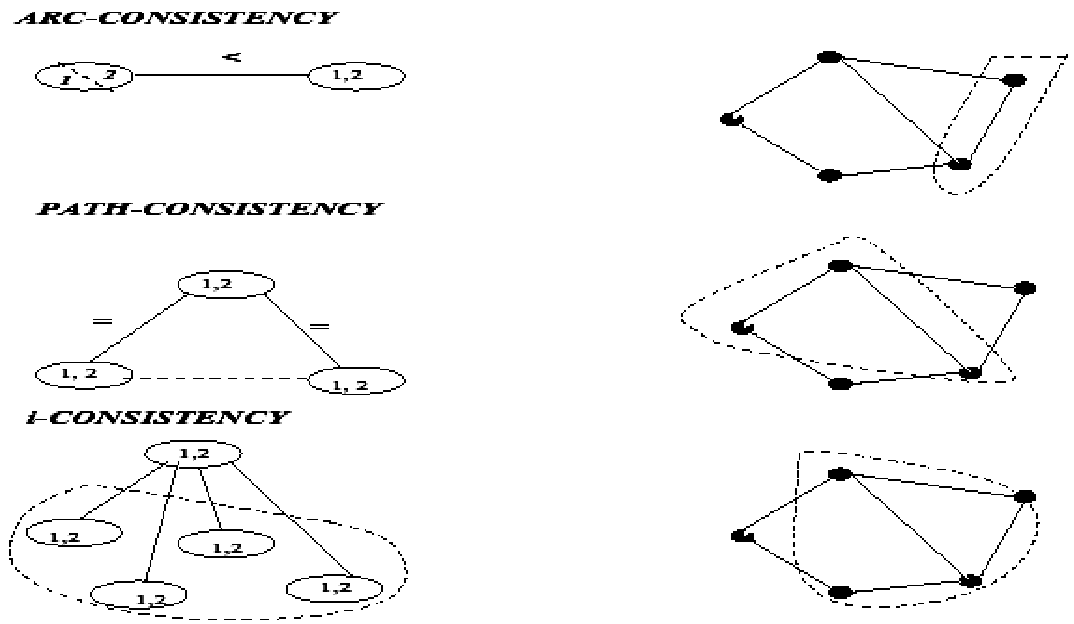


Figure 3.17: The scope of consistency enforcing: (a) arc-consistency, (b) path-consistency, (c) i-consistency

Higher levels of consistency, global-consistency

Definition 3.4.1 (i-consistency, global consistency) *Given a general network of constraints $\mathcal{R} = (X, D, C)$, a relation $R_S \in C$ where $|S| = i - 1$ is i-consistent relative to a variable y not in S iff for every $t \in R_S$, there exists a value $a \in D_y$, s.t. (t, a) is consistent. A network is* i-consistent *iff given any consistent instantiation of any $i - 1$ distinct variables, there exists an instantiation of any ith variable such that the i values taken together satisfy all of the constraints among the i variables. A network is* strongly i-consistent *iff it is j-consistent for all $j \leq i$. A strongly n-consistent network, where n is the number of variables in the network, is called* globally consistent.

Revise-i

REVISE-$i(\{x_1, x_2, ..., x_{i-1}\}, x_i)$

input: a network $\mathcal{R} = (X, D, C)$

output: a constraint R_S, $S = \{x_1, ..., x_{i-1}\}$ i-consistent relative to x_i.

1. **for** each instantiation $\bar{a}_{i-1} = (< x_1, a_1 >, < x_2, a_2 >, ..., < x_{i-1}, a_{i-1} >)$ do,
2. **if** no value of $a_i \in D_i$ exists s.t. $(\bar{a}_{i-1}, a_i)$ is consistent

   then delete $\bar{a}_{i-1}$ from R_S

   (Alternatively, let $\mathcal{S}$ be the set of all subsets of $\{x_1, ..., x_i\}$ that contain x_i and appear as scopes of constraints of $\mathcal{R}$, then

   $R_S \leftarrow R_S \cap \pi_S(\bowtie_{S' \subseteq \mathcal{S}} R_{S'}))$
3. **endfor**

Figure 3.14: Revise-i

- Complexity: for binary constraints $O(k^i)$
- For arbitrary constraints: $O((2k)^i)$

4-queen example

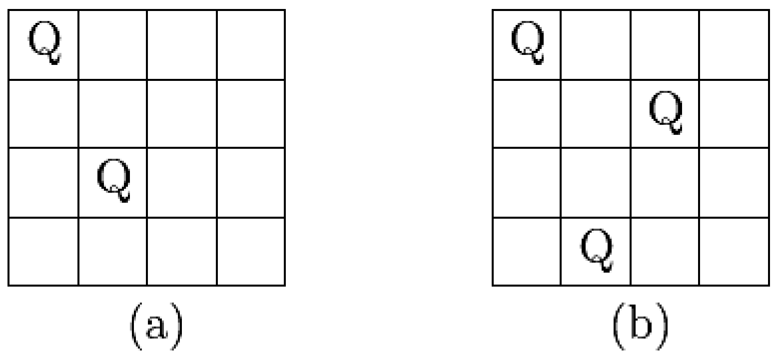


Figure 3.13: (a) Not 3-consistent; (b) Not 4-consistent

I-consistency

I-CONSISTENCY($\mathcal{R}$)

input: a network $\mathcal{R}$.

output: an i-consistent network equivalent to $\mathcal{R}$.

1. **repeat**
2. **for** every subset $S \subseteq X$ of size $i-1$, and for every x_i, do
3. let $\mathcal{S}$ be the set of all subsets in of $\{x_1, ..., x_i\}$ $scheme(\mathcal{R})$ that contain x_i
4. $R_S \leftarrow R_S \cap \pi_S(\bowtie_{S' \in \mathcal{S}} R_{S'})$ (this is Revise-i(S, x_i))
6. **endfor**
7. **until** no constraint is changed.

Figure 3.15: i-consistency-1

Theorem 3.4.3 (complexity of i-consistency) *The time and space complexity of brute-force i-consistency $O(2^i(nk)^{2i})$ and $O(n^i k^i)$, respectively. A lower bound for enforcing i-consistency is $\Omega(n^i k^i)$.* □

Arc-consistency for non-binary constraints:
Generalized arc-consistency

Definition 3.5.1 (generalized arc-consistency) *Given a constraint network* $\mathcal{R} = (\mathcal{X}, \mathcal{D}, \mathcal{C})$, *with* $R_S \in C$, *a variable* x *is* arc-consistent *relative to* R_S *if and only if for every value* $a \in D_x$ *there exists a tuple* $t \in R_S$ *such that* $t[x] = a$. t *can be called a support for* a. *The constraint* R_S *is called arc-consistent iff it is arc-consistent relative to each of the variables in its scope and a constraint network is arc-consistent if all its constraints are arc-consistent.*

$$D_x \leftarrow D_x \cap \pi_x (R_S \otimes D_{S-\{x\}})$$

Complexity: O(t k), t bounds number of tuples.

Relational arc-consistency:

$$R_{S-\{x\}} \leftarrow \pi_{S-\{x\}} (R_S \otimes D_x)$$

Examples of generalized arc-consistency

- {x+y+z <= 15, z >= 13} → x<=2, y<=2

- Example of relational arc-consistency

$$\{A \wedge B \rightarrow G, \neg G\}, \Rightarrow \neg A \vee \neg B$$

More arc-based consistency

- Global constraints: e.g., all-different constraints
  - Special semantic constraints that appears often in practice and a specialized constraint propagation. Used in constraint programming.
- Bounds-consistency: pruning the boundaries of domains
- Do exercise 16

Example for alldiff

- A = {3,4,5,6}
- B = {3,4}
- C= {2,3,4,5}
- D= {2,3,4}
- E = {3,4}
- Alldiff (A,B,C,D,E}
- Arc-consistency does nothing
- Apply GAC to sol(A,B,C,D,E)?
- → A = {6}, F = {1}….
- Alg: bipartite matching kn^1.5
- (Lopez-Ortiz, et. Al, IJCAI-03 pp 245 (A fast and simple algorithm for bounds consistency of alldifferent constraint)

Global constraints

- Alldifferent
- Sum constraint
- Global cardinality constraint (a value can be assigned a bounded number of times)
- The cummulative constraint (related to scheduling tasks)

Bounds consistency

Definition 3.5.4 (bounds consistency) *Given a constraint C over a scope S and domain constraints, a variable $x \in S$ is bounds-consistent relative to C if the value $min\{D_x\}$ (respectively, $max\{D_x\}$) can be extended to a full tuple t of C. We say that t supports $min\{D_x\}$. A constraint C is* bounds-consistent *if each of its variables is bounds-consistent.*

Bounds consistency for Alldifferent constraints

Example 3.5.5 Consider the constraint problem with variables $x_1, ...x_6$, each with domains $1, ..., 6$, and constraints:

$$C_1 : x_4 \geq x_1 + 3, \quad C_2 : x_4 \geq x_2 + 3, \quad C_3 : x_5 \geq x_3 + 3, \quad C_4 : x_5 \geq x_4 + 1,$$

$$C_5 : alldifferent\{x_1, x_2, x_3, x_4, x_5\}$$

The constraints are not bounds consistent. For example, the minimum value 1 in the domain of x_4 does not have support in constraint C_1 as there is no corresponding value for x_1 that satisfies the constraint. Enforcing bounds consistency using constraints C_1 through C_4 reduces the domains of the variables as follows: $D_1 = \{1, 2\}$, $D_2 = \{1, 2\}$, $D_3 = \{1, 2, 3\}$ $D_4 = \{4, 5\}$ and $D_5 = \{5, 6\}$. Subsequently, enforcing bounds consistency using constraints C_5 further reduces the domain of C to $D_3 = \{3\}$.Now constraint C_3 is no longer bound consistent. Reestablishing bounds consistency causes the domain of x_5 to be reduced to $\{6\}$. Is the resulting problem already arc-consistent? □

Boolean constraint propagation

- (A V ~B) and (B)
  - B is arc-consistent relative to A but not vice-versa
- Arc-consistency achieved by resolution:

  res((A V ~B),B) = A

Given also (B V C), path-consistency means:

res((A V ~B),(B V C)) = (A V C)

What can generalized arc-consistency do to cnfs?

Relational arc-consistency rule = unit-resolution

Boolean constraint propagation

Example: party problem

- If Alex goes, then Becky goes: $\mathbf{A} \rightarrow \mathbf{B}$ $(\text{or}, \neg\mathbf{A} \vee \mathbf{B})$
- If Chris goes, then Alex goes: $\mathbf{C} \rightarrow \mathbf{A}$ $(\text{or}, \neg\mathbf{C} \vee \mathbf{A})$
- **Query:**

  Is it possible that Chris goes to the party but Becky does not?

$$\Updownarrow$$

$$\text{Is } \textit{propositional theory}$$
$$\varphi = \{\neg\mathbf{A} \vee \mathbf{B},\ \neg\mathbf{C} \vee \mathbf{A},\ \neg\mathbf{B},\ \mathbf{C}\} \text{ satisfiable?}$$

Constraint propagation for Boolean constraints: Unit propagation

Procedure UNIT-PROPAGATION
Input: A cnf theory, φ, $d = Q_1, ..., Q_n$.
Output: An equivalent theory such that every unit clause does not appear in any non-unit clause.

1. queue = all unit clauses.
2. **while** queue is not empty, do.
3. $T \leftarrow$ next unit clause from Queue.
4. **for** every clause β containing T or $\neg T$
5. **if** β contains T delete β (subsumption elimination)
6. **else**, For each clause $\gamma = resolve(\beta, T)$.
   if γ, the resolvent, is empty, the theory is unsatisfiable.
7. **else**, add the resolvent γ to the theory and delete β.
   if γ is a unit clause, add to Queue.
8. **endfor**.
9. **endwhile**.

Theorem 3.6.1 *Algorithm* UNIT-PROPAGATION *has a linear time complexity.*

Algorithms for relational and generalized arc-cnsistency

- Think about the following:
  - GAC-i apply AC-i to the dual problem when singleton variables are explicit: the bi-partite representation.
  - What is the complexity?
  - Relational arc-consistency: imitate unit propagation.
  - Apply AC-1 on the dual problem where each subset of a scope is presented.
  - Is unit propagation equivalent to AC-4?

Consistency for numeric constraints

$$x \in [1,10],\ y \in [5,15],$$

$$x + y = 10$$

$$arc - consistency \Rightarrow x \in [1,5],\ y \in [5,9]$$

$$by - adding - x + y = 10, -y \le -5$$

$$z \in [-10,10],$$

$$y + z \le 3$$

$$path - consistency \Rightarrow x - z \ge 7$$

$$obtained - by - adding, x + y = 10, -y - z \ge -3$$

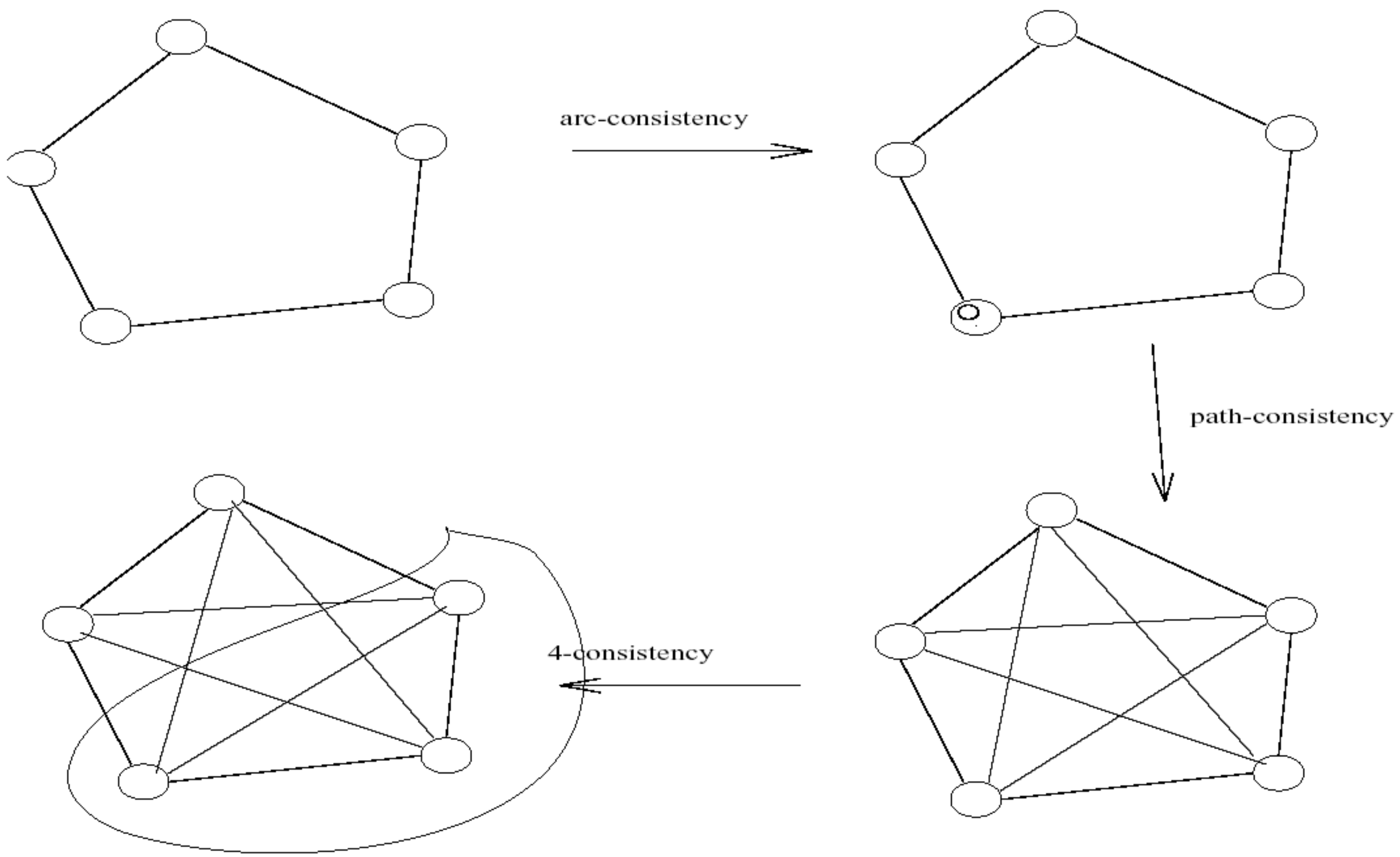
Tractable classes

Theorem 3.7.1 *1. The consistency binary constraint networks having no cycles can be decided by arc-consistent*

2. The consistency of binary constraint networks with bi-valued domains can be decided by path-consistency,

3. The consistency of Horn cnf theories can be decided by unit propagation.

Changes in the network graph as a result of arc-consistency, path-consistency and 4-consistency.



Distributed arc-consistency (Constraint propagation)

- **Implement AC-1 distributedly.**

$$D_i \leftarrow D_i \cap \pi_i(R_{ij} \otimes D_j)$$

- **Node x_j sends the message to node x_i**

$$h_i^j \leftarrow \pi_i(R_{ij} \otimes D_j)$$

$$D_i \leftarrow D_i \cap h_i^j$$

- **Node x_i updates its domain:**

- **Generalized arc-consistency can be implemented distributedly: sending messages between constraints over the dual graph:**

$$R_{S-\{x\}} \leftarrow \pi_{S-\{x\}}(R_S \otimes D_x)$$

Distributed relational arc-consistency example

The message that R2 sends to R1 is

R1 updates its relation and domains and sends messages to neighbors

R_1

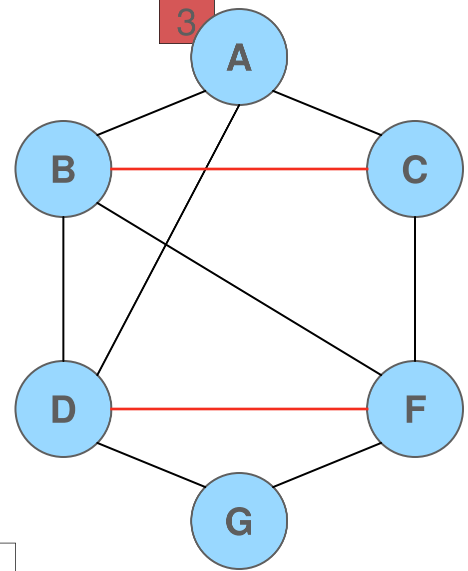
A
1
2
3

R_2

A	B
1	2
1	3
2	1
2	3
3	1
3	2

R_3

A	C
1	2
3	2



R_4

A	B	D
1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

R_5

B	C	F
1	2	3
3	2	1

R_6

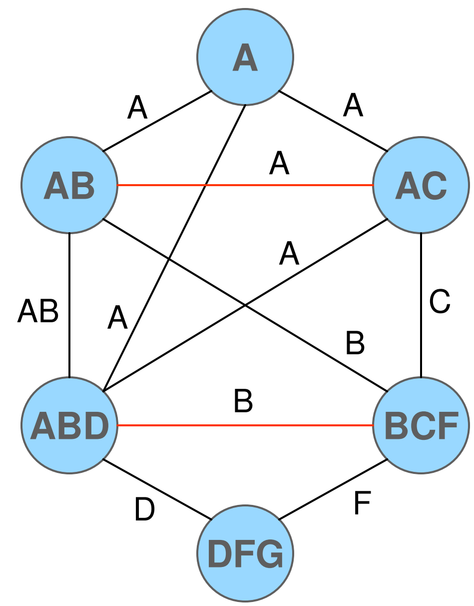
D	F	G
1	2	3
2	1	3

Distributed Arc-Consistency

DR-AC can be applied to the dual problem of any constraint network.

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \quad (1)$$

$$D_i \leftarrow D_i \cap (\bowtie_{k \in ne(i)} D_k^i) \quad (2)$$



b) Constraint network

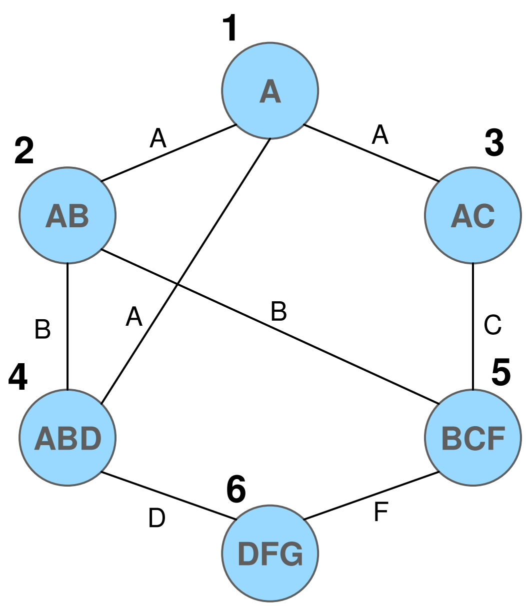
DR-AC on a dual join-graph

R_1

A
1
2
3

R_2

A	B
1	2
1	3
2	1
2	3
3	1
3	2



R_3

A	C
1	2
3	2

R_4

A	B	D
1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

R_5

B	C	F
1	2	3
3	2	1

R_6

D	F	G
1	2	3
2	1	3

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \qquad (1)$$

Iteration 1

R_1

A
1
2
3

h_2^1

A
1
2
3

h_3^1

A
1
3

h_4^1

A
1
2
3

h_5^2

B
1
3

h_4^2

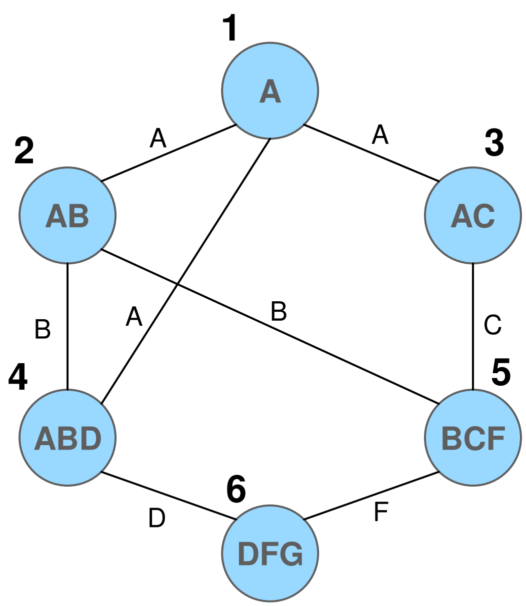
B
1
2
3

h_1^2

A
1
2
3

R_2

A	B
1	2
1	3
2	1
2	3
3	1
3	2



R_3

A	C
1	2
3	2

h_1^3

A
1
2
3

h_5^3

C
2

h_6^4

D
1
2

h_2^4

B
1
2
3

h_1^4

A
1
2
3

R_4

A	B	D
1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

R_5

B	C	F
1	2	3
3	2	1

h_2^5

B
1
2
3

h_3^5

C
2

h_6^5

F
1
3

R_6

D	F	G
1	2	3
2	1	3

h_4^6

D
1
2
3

h_5^6

F
1
3

$$R_i \leftarrow R_i \cap (\bowtie_{k \in ne(i)} h_k^i) \quad (2)$$

Iteration 1

R_1

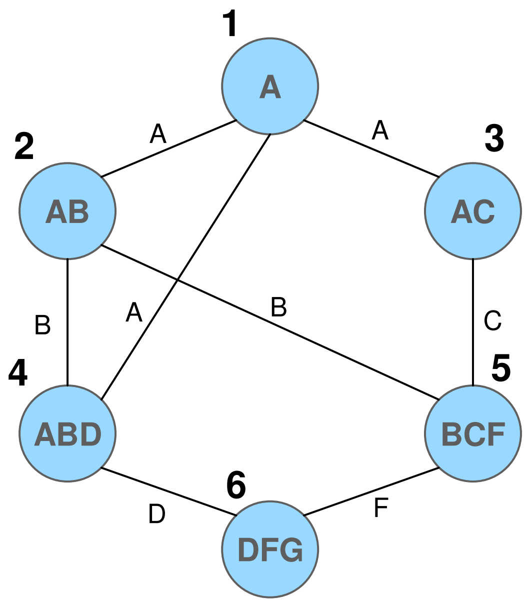
A
1
3

R_2

A	B
1	3
2	1
2	3
3	1

R_3

A	C
1	2
3	2



R_4

A	B	D
1	3	2
2	3	1
3	1	2
3	2	1

R_5

B	C	F
1	2	3
3	2	1

R_6

D	F	G
2	1	3

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \qquad (1)$$

Iteration 2

R_1

A
1
3

h_2^1

A
1
2
3

h_3^1

A
1
3

h_4^1

A
1
2
3

h_5^2

B
1
3

h_4^2

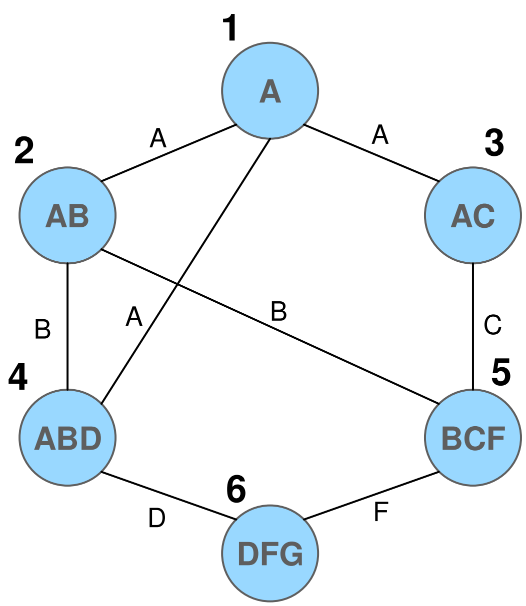
B
1
2
3

h_1^2

A
1
3

R_2

A	B
1	3
2	1
2	3
3	1



R_3

A	C
1	2
3	2

h_1^3

A
1
3

h_5^3

C
2

h_6^4

D
2

h_2^4

B
1
3

h_1^4

A
1
3

R_4

A	B	D
1	3	2
2	3	1
3	1	2
3	2	1

R_5

B	C	F
1	2	3
3	2	1

h_2^5

B
1
3

h_3^5

C
2

h_6^5

F
1

R_6

D	F	G
2	1	3

h_4^6

D
1
2

h_5^6

F
1
3

$$R_i \leftarrow R_i \cap (\bowtie_{k \in ne(i)} h_k^i) \quad (2)$$

Iteration 2

R_1

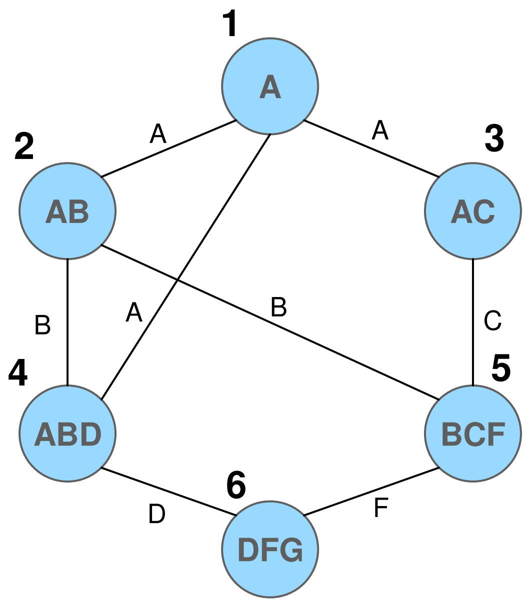
A
1
3

R_2

A	B
1	3
3	1

R_3

A	C
1	2
3	2



R_4

A	B	D
1	3	2
3	1	2

R_5

B	C	F
3	2	1

R_6

D	F	G
2	1	3

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \qquad (1)$$

Iteration 3

R_1

A
1
3

h_2^1

A
1
3

h_3^1

A
1
3

h_4^1

A
1
3

h_5^2

B
3

h_4^2

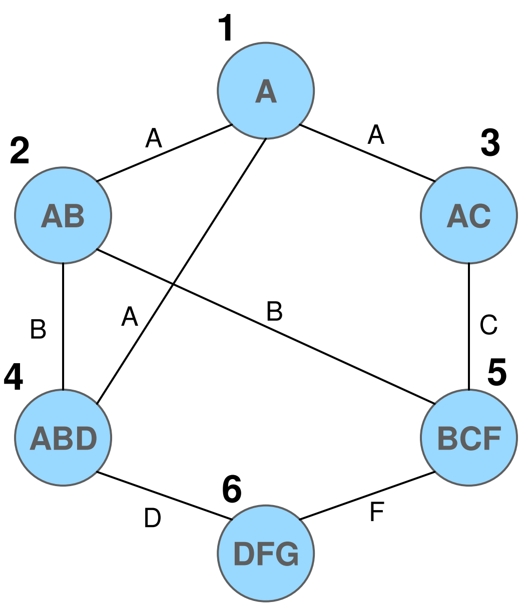
B
1
3

h_1^2

A
1
3

R_2

A	B
1	3
3	1



R_3

A	C
1	2
3	2

h_1^3

A
1
3

h_5^3

C
2

h_6^4

D
2

h_2^4

B
1
3

h_1^4

A
1
3

R_4

A	B	D
1	3	2
3	1	2

R_5

B	C	F
3	2	1

h_2^5

B
1
3

h_3^5

C
2

h_6^5

F
1

R_6

D	F	G
2	1	3

h_4^6

D
2

h_5^6

F
1

$$R_i \leftarrow R_i \cap (\bowtie_{k \in ne(i)} h_k^i) \qquad (2)$$

Iteration 3

R_1

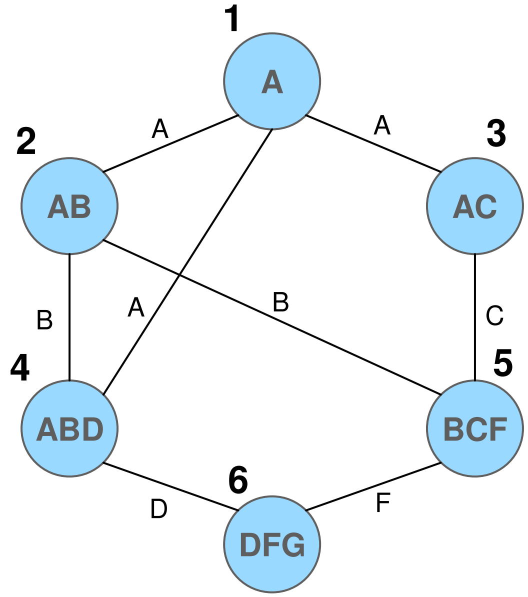
A
1
3

R_2

A	B
1	3

R_3

A	C
1	2
3	2



R_4

A	B	D
1	3	2
3	1	2

R_5

B	C	F
3	2	1

R_6

D	F	G
2	1	3

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \qquad (1)$$

Iteration 4

R_1

A
1
3

h_2^1

A
1

h_3^1

A
1
3

h_4^1

A
1
3

h_5^2

B
3

h_4^2

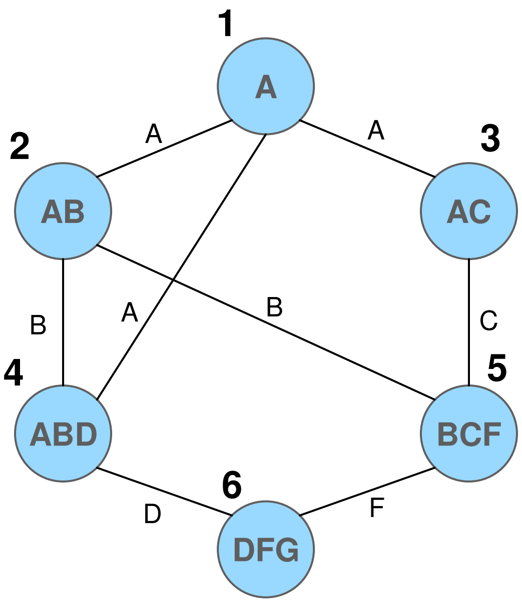
B
1
3

h_1^2

A
1
3

R_2

A	B
1	3



R_3

A	C
1	2
3	2

h_1^3

A
1
3

h_5^3

C
2

h_6^4

D
2

h_2^4

B
1
3

h_1^4

A
1
3

R_4

A	B	D
1	3	2
3	1	2

R_5

B	C	F
3	2	1

h_2^5

B
3

h_3^5

C
2

h_6^5

F
1

R_6

D	F	G
2	1	3

h_4^6

D
2

h_5^6

F
1

$$R_i \leftarrow R_i \cap (\bowtie_{k \in ne(i)} h_k^i) \quad (2)$$

Iteration 4

R_1

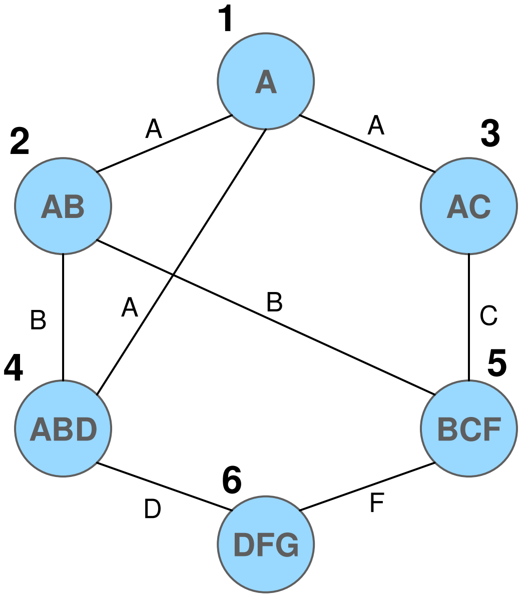
A
1

R_2

A	B
1	3

R_3

A	C
1	2
3	2



R_4

A	B	D
1	3	2

R_5

B	C	F
3	2	1

R_6

D	F	G
2	1	3

$$h_i^j \leftarrow \pi_{l_{ij}}(R_i \bowtie (\bowtie_{k \in ne(i)} h_k^i)) \quad (1)$$

Iteration 5

R_1

A
1

h_2^1

A
1

h_3^1

A
1

h_4^1

A
1

h_5^2

B
3

h_4^2

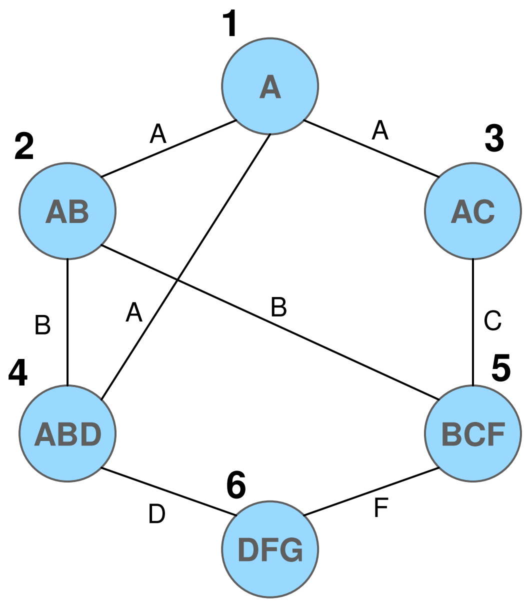
B
3

h_1^2

A
1

R_2

A	B
1	3



R_3

A	C
1	2
3	2

h_1^3

A
1

h_5^3

C
2

h_6^4

D
2

h_2^4

B
3

h_1^4

A
1

R_4

A	B	D
1	3	2

R_5

B	C	F
3	2	1

h_2^5

B
3

h_3^5

C
2

h_6^5

F
1

R_6

D	F	G
2	1	3

h_4^6

D
2

h_5^6

F
1

$$R_i \leftarrow R_i \cap (\bowtie_{k \in ne(i)} h_k^i) \quad (2)$$

Iteration 5

R_1

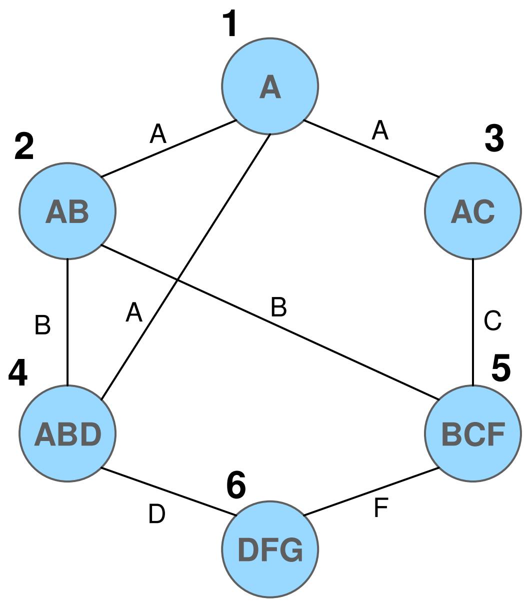
A
1

R_2

A	B
1	3

R_3

A	C
1	2



R_4

A	B	D
1	3	2

R_5

B	C	F
3	2	1

R_6

D	F	G
2	1	3